

## Ed meme recap:

## don't write bugs.

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Jocelyn Zhang staff
Jocelyn Zhang staff
16 hours ago in Resources
On endorsed
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## Questions on lecture content? Or about cats?

## Stress

- 429 H is not an easy class
- Lots of new materials
- Unfamiliar programming environments
- Fast, often relentless pace
- Struggling in this course is normal
- There will be times you won't know the answer of the solution
- This is expected-we want we everyone to succeed, but the only way we can help is if you ask for it
- If you find yourself overly overwhelmed or spending more time on this class than you think you should be, please reach out to Dr. Gheith or the TAs
- We can help out as far as the class goes
- We can provide other resources where we are not able to help

Mental health resource available at UT

P5

## Poll

How's your status on P5?
A. What's P5?
B. I've heard of it
C. I've cloned the starter code and/or looked through it
D. I've started planning/writing code
E. I'm mostly done but might still have bugs
F. P5 any\% speedrun

## PSA, run your code with Gheith's GCC version

We will use this version with -O3 to grade.

The instructions:
https://edstem.org/us/courses/53774/discussion/4470800

P6

## First, what's a subroutine?

- Is it the same thing as a function?
- Why is the name subroutine significant?
- How do we keep track of the execution state when going between different subroutines?


## What's a Coroutine? Something ducks walk on?

- Consider an ordinary function
- May make calls to subroutines, but always executes sequentially
- Can only run one of these at a time, given 1 CPU
- What if we want to run a lot of functions, interleaving their execution?
- Need a notion of a "suspendable function"---a function that we can stop and resume
- This way, we can make the PC jump between routines
- Concurrent but not parallel
- How do we decide when to switch routines?
- Preemptive - somehow force the CPU to switch tasks
- Cooperative - define specific "yield points" at which the CPU switches tasks


## What's a Coroutine? Something ducks walk on?

- Consider an ordinary function
- Uses function parameters and return values to exchange information
- What if we want separate coroutines to communicate?
- Need an explicit mechanism (ene that says bad words alot)
- Useful abstraction: channels
- Channels
- Object passed to one or more routines
- Interface consists of send and receive functions
- What should these do?
- What data structure does this sound like?


## Fun Activity Time!

- Groups of 2 will pretend to be coroutines.
- One coroutine gets to execute at a time.
- Coroutines will send each other messages. A send cannot finish unless the receiver is also trying to receive at the same time. In this case, the send will block the sender from executing. Ditto for receives.
- When a coroutine yields, any unblocked coroutine gets the chance to run. The TAs will use a tiebreaker to see who gets to run, and the rest of the class also helps in this tiebreaker.
- For the purposes of this activity, sends and receives will always yield.
- The coroutines only get to look at their own instructions, everyone else can look at the entire program. It's posted on Ed now!


## Coroutines in Practice - Saving Progress

- Consider an ordinary function
- All necessary information for execution and returning are saved on the stack and registers.
- With more than one routine, we must save this information
- How do we keep track of every routine's state?


## What is a stack?



## What is a stack in computers?



## What are stacks good for?

- Storing temporary/local variables
- Storing return addresses for function calls
- Storing function arguments


## More on stacks

- Where is the stack usually?
- How do you find it?
- Does it have to be this way?


## P6 - Some Assembly Required

- Context switching via magic.S
- takes the active (running) coroutine and swaps it for a different coroutine, saving the first one's state after deactivating it
- general flow of a context-switching function:
- Record the complete state of the currently running coroutine and write it (implicitly or explicitly) into a Routine (CCB) struct
- Read the state of the next coroutine from its Routine struct, and restore it
- return (but to where????)
- state in the Routine struct can be saved explicitly (through struct fields) or implicitly (by pushing to the stack and saving only the stack pointer in the Routine struct


## Poll

## Which ISA will P6 use?

1. ARM
2. $x 86$
3. RISC-V
4. PowerPC
5. MIPS


## P6 - Some Assembly Required

- Calling Convention:
- X86:
- Args go in \%rdi, \%rsi, \%rdx, \%rcx for integral types (in that order).
- You will not need more than that many arguments
- Return value goes in \%rax
- Callee saved registers: \%rsp, \%rbp, \%rbx, \%r12, \%r13,\%r14, \%r15
- ARM:
- Argos go in $x 0, x 1, \times 2, \times 3$ for integral types (in that order).
- You will not need more than that many arguments
- Return value goes in $x 0$
- Callee saved registers: $\times 19-\times 28$


## P6 - Some Assembly Required

- Where might you need to return a value from assembly?
- Which registers should you save? Caller/callee/both?


## P6 - Some Assembly Required

- how can we tell the compiler that there is an asm function called magic?

```
o extern <return_type> magic(<args>...);
```

- what do we do in assembly to create the function definition?
- .global magic
o magic:
o ret
- similar method to achieve the opposite effect
- .extern c_function
- call c_function / bl c_function


## Coroutines in Practice - The Real World

- "lazily evaluated" computations
- Generators in Python are an example of this
- simpler iterator implementations
- Remember writing an iterator for Boggle in 314H? How could it be easier with coroutines?
- asynchronous programming
- When you have to wait for the network, disk, etc., but want to keep doing things in the meantime, coroutines can make that really nice and simple


## Quiz,everyone say YIPPEE! REVIEW

## Question 1

[1 point] The concept of settling time is how long a logical gate takes to converge to an output after receiving an input. After a gate "settles", we can reliably read the output of the gate. If each NOT gate takes 5 ns to settle and each AND takes 7 ns to settle, evaluate the time it takes for the circuit below to settle on its output values.


## Question 2a

## creating gates with nor

not:


## Question 2b

## creating gates with nor

and:


## Question 2c

## creating gates with nor

xor:


## Question 3

[ 0.5 points] What is your favorite instruction in ARM? Briefly explain.

- If you were creating a rubric for this question, what would you do?
- What would be required for full credit?


## Question 3

[ 0.5 points] What is your favorite instruction in ARM? Briefly explain.

- If you were creating a rubric for this question, what would you do?
- What would be required for full credit?
" $B$ is my favorite instruction because it is fast to type with just one character"
- Y'all did well on this


## Question 4

[1 point] The CSETEQ instruction on ARM will set a register to 1 if the zero flag is set (the previous comparison resulted in equality), or 0 otherwise. Implement the functionality of this instruction using only (conditional or unconditional) branches and move statements: fill in the following code and produce the output in the register $x 7$.

## Example of CSETEQ:

// Assume x0 = 3, x1 = 4, x2 = 4
cmp x1, x0
cseteq x3 // x3 is now 0
cmp x1, x2
cseteq $x 4$ // $x 4$ is now 1

## Answer space:

cmp x5, x6
// implement cseteq x7 here

## Question 4

[1 point] The CSETEQ instruction on ARM will set a register to 1 if the zero flag is set (the previous comparison resulted in equality), or 0 otherwise. Implement the functionality of this instruction using only (conditional or unconditional) branches and move statements: fill in the following code and produce the output in the register $x 7$.
cmp x5, x6
b.eq equals
mov x7, \#0
b end
equals:
mov x1, \#1
end:

```
if (x5 != x6) {
        <7 = 0;
} else {
    x7 = 1;
}
```


## Question 4

[1 point] The CSETEQ instruction on ARM will set a register to 1 if the zero flag is set (the previous comparison resulted in equality), or 0 otherwise. Implement the functionality of this instruction using only (conditional or unconditional) branches and move statements: fill in the following code and produce the output in the register $x 7$.

```
cmp x5, x6
mov x7, #0
b.ne end
mov x1, #1
end:
```

```
x7 = 0;
if (x5 == x6) {
    x7 = 1;
}
```


## Question 5

## bitstring $b_{-1} b_{-2} \ldots b_{-(N-1)} b_{-N}$

## value $\mathrm{b}_{-1} * 2^{-1}+\mathrm{b}_{-2} * 2^{-2}+\ldots+\mathrm{b}_{-(\mathrm{N}-1)} * 2^{-\left(\mathrm{N}_{-1}\right)}+\mathrm{b}_{-\mathrm{*}} * 2^{-N}$.

- Basic addition works the same as with regular binary numbers, adding bits from right to left and propagating carry values
- If you have N bits, numbers that are not of the form $\mathrm{k} / 2^{\wedge} \mathrm{N}$ where k is some integer cannot be represented by our system


## Question 6

Two's Complement

- Arithmetic is simple
- Adding and subtracting are the same
- $\quad 2^{N}$ distinct values representable
- Negation is hard
- Sign extension is hard

Linner Linner Chicken Dinner

- Arithmetic is hard
- Adding and subtracting are not the same
- $2^{\mathrm{N}}-1$ distinct values representable
- Negation is easy
- Sign extension is trivial


## Question 6

Two's Complement

- Arithmetic is simple
- Adding and subtracting are the same
- $\quad 2^{\mathrm{N}}$ distinct values representable
- Negation is hard
- Sign extension is hard

One's Complement

- Arithmetic is slightly harder
- Adding and subtracting are nearly the same
- $2^{\mathrm{N}}-1$ distinct values representable
- Negation is slightly less hard
- Sign extension slightly less hard


## Question 6 (Bonus)

- $a+b+c a r r y$ bit
- Really, it's as shrimple as that

Questions?

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